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# ON HEILBRONN'S PROBLEM IN HIGHER DIMENSION

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Heilbronn conjectured that given arbitrary n points in the 2-dimensional unit square  $[0,1]^2$ , there must be three points which form a triangle of area at most  $O(1/n^2)$ . This conjecture was disproved by a nonconstructive argument of Komlós, Pintz and Szemerédi [10] who showed that for every n there is a configuration of n points in the unit square  $[0,1]^2$  where all triangles have area at least  $\Omega(\log n/n^2)$ . Considering a generalization of this problem to dimensions  $d \geq 3$ , Barequet [3] showed for every n the existence of n points in the d-dimensional unit cube  $[0,1]^d$  such that the minimum volume of every simplex spanned by any (d+1) of these n points is at least  $\Omega(1/n^d)$ . We improve on this lower bound by a logarithmic factor  $\Theta(\log n)$ .

## 1. Introduction

An old conjecture of Heilbronn states that for every distribution of n points in the 2-dimensional unit square  $[0,1]^2$  (or unit disc) there are three distinct points which form a triangle of area at most  $c/n^2$  for some constant c>0. Erdős observed that this conjecture, if true, would be best possible, as, for n a prime, the points  $(i,i^2 \mod n)_{i=0,\dots,n-1}$  in the  $n\times n$  grid would show after rescaling, see [2]. However, Komlós, Pintz and Szemerédi [10] disproved Heilbronn's conjecture by showing for every n the existence of a configuration of n points in  $[0,1]^2$  with every three of these n points forming a triangle of area at least  $c' \cdot \log n/n^2$  for some constant c' > 0. This existence argument was made constructive in [5], where a deterministic polynomial time algorithm was given, which finds n points in  $[0,1]^2$  achieving this lower bound

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 $\Omega(\log n/n^2)$  on the minimum triangle area. Upper bounds on Heilbronn's triangle problem were given by Roth in a series [12–16] of papers and by Schmidt [18], see Rothschild and Straus [17] for related results, and the currently best upper bound  $O(n^{-8/7+\varepsilon})$  for every fixed  $\varepsilon > 0$  is due to Komlós, Pintz and Szemerédi [9].

Recently, Barequet [3] considered a d-dimensional version of Heilbronn's problem. For given (d+1) vectors  $p_1, \ldots, p_{d+1} \in \mathbb{R}^d$  the set  $\{\sum_{i=1}^{d+1} \lambda_i \cdot \}$  $p_i \mid \sum_{i=1}^{d+1} \lambda_i = 1; \ \lambda_1, \dots, \lambda_{d+1} \geq 0$  is called a *simplex*. For fixed dimension sion  $d \geq 3$ , Barequet showed for every n, that there exist n points in the d-dimensional unit cube  $[0,1]^d$  such that the minimum volume of every simplex spanned by any (d+1) of these n points is at least  $\Omega(1/n^d)$ . He gave three different approaches towards a solution of the problem. The first one, for dimension d=3, uses a Greedy-type argument, i.e., adding to given points a new point as long as possible, such that no two points are too close, no three points form a triangle of too small area and no four points form a tetrahedron of too small volume (see also [18] for the case d=2). With this he obtained a configuration of n points in the 3-dimensional unit cube  $[0,1]^3$ such that the minimum volume of every tetrahedron is at least  $\Omega(1/n^4)$ . The second approach, which yields a better lower bound, was worked out for every fixed dimension d > 3 and uses a random argument: if 2n points are dropped uniformly at random and independently of each other in the d-dimensional unit cube  $[0,1]^d$ , then the expected number of simplices with volume at most  $c_d/n^d$  is at most n, where  $c_d > 0$  is a constant. Deleting one point from every such small simplex yields the existence of n points in  $[0,1]^d$  with every simplex formed by (d+1) of these points having volume at least  $\Omega(1/n^d)$ . The third approach however is similar to Erdős one's (and according to Bollobás [6] was known to him) and is an explicit construction, namely taking the points  $P_k = 1/n \cdot (k^j \mod n)_{j=1,\dots,d}$  for  $k=0,1,\dots,n-1$  on the moment curve. The volume of every simplex is given by the determinant of a Vandermonde matrix, which is not equal to 0 for n a prime, multiplied by  $\Theta(1/n^d)$  and this gives minimum value at least  $\Omega(1/n^d)$ .

Note, that the corresponding problem in dimension d=1 is trivial as n equidistant points in the unit interval [0,1] show.

Here we will improve Barequet's lower bound for dimensions  $d \ge 3$ , using a probabilistic existence argument, by a factor  $\Theta(\log n)$ :

**Theorem 1.1.** For every fixed integer  $d \ge 2$  and for every n there exists a configuration of n points in the unit cube  $[0,1]^d$  such that the volume of every simplex spanned by any (d+1) of these points is at least  $\Omega(\log n/n^d)$ .

## 2. Simplices with Small Volume and Hypergraphs

In our arguments we will use hypergraphs. The parameters *independence* number of a hypergraph and 2-cycles will be important in our considerations:

**Definition 2.1.** Let  $\mathcal{G} = (V, \mathcal{E})$  be a hypergraph with vertex set V and edge set  $\mathcal{E}$  where each edge  $E \in \mathcal{E}$  satisfies  $E \subseteq V$ . A hypergraph  $\mathcal{G} = (V, \mathcal{E})$  is k-uniform if every edge  $E \in \mathcal{E}$  contains exactly k vertices.

A subset  $I \subseteq V$  is called *independent* if I contains no edge  $E \in \mathcal{E}$ . The largest size of an independent set in  $\mathcal{G}$  is called the *independence number*  $\alpha(\mathcal{G})$ .

In a k-uniform hypergraph  $\mathcal{G} = (V, \mathcal{E})$ ,  $k \geq 3$ , a 2-cycle is a pair  $\{E_1, E_2\}$  of distinct edges  $E_1, E_2 \in \mathcal{E}$  with  $|E_1 \cap E_2| \geq 2$ . A 2-cycle  $\{E_1, E_2\}$  in  $\mathcal{G}$  is called (2, j)-cycle if  $|E_1 \cap E_2| = j$ , where j = 2, ..., k - 1.

We will reformulate the geometrical problem considered by Barequet as a problem of finding in an appropriately defined hypergraph a large independent set. For a given set  $S \subseteq [0,1]^d$  of points we form a (d+1)-uniform hypergraph with vertex set being this set S of points in  $[0,1]^d$ . The edges are determined by all subsets of (d+1) points from S, which form a simplex of 'small' volume, to be specified later. An independent set in this hypergraph corresponds to a set of points in  $[0,1]^d$ , where no simplex has 'small' volume. In order to show the existence of a large independent set, we will use the following result of Ajtai, Komlós, Pintz, Spencer and Szemerédi [1], stated here in a variant proven in [7]:

**Theorem 2.2** ([1],[7]). Let  $k \geq 3$  be a fixed integer. Let  $\mathcal{G} = (V, \mathcal{E})$  be a k-uniform hypergraph on |V| = n vertices and with average degree  $t^{k-1} = k \cdot |\mathcal{E}|/n$ . If  $\mathcal{G}$  does not contain any 2-cycles, then the independence number  $\alpha(\mathcal{G})$  satisfies for some constant  $c_k > 0$ :

$$\alpha(\mathcal{G}) \ge c_k \cdot \frac{n}{t} \cdot (\log t)^{\frac{1}{k-1}}$$
.

In recent years, several applications and also an algorithmic version of Theorem 2.2 have been found, compare [4]. Here we will give a another application of this deep result.

In d dimensions the volume of a simplex determined by the points  $P_1, \ldots, P_{d+1} \in [0,1]^d$  is given by  $\operatorname{vol}(P_1, \ldots, P_{d+1}) := 1/d \cdot G \cdot h$ , where G is the volume of the simplex determined by the points  $P_1, \ldots, P_d$  (in the corresponding (d-1)-dimensional subspace) and h is the Euclidean distance of the point  $P_{d+1}$  from the hyperplane given by  $P_1, \ldots, P_d$ . Thus, if  $h_k$  denotes the

Euclidean distance of  $P_k$  from the hyperplane determined by  $P_1, \ldots, P_{k-1}, k=2,\ldots,d+1$ , then

$$vol(P_1, ..., P_{d+1}) = \frac{1}{d!} \cdot \prod_{k=2}^{d+1} h_k.$$

In the following we will prove Theorem 1.1.

**Proof.** In the *d*-dimensional unit cube  $[0,1]^d$  we drop  $n^{1+\varepsilon}$  points uniformly at random and independently of each other, where  $\varepsilon$  is a small constant with  $0 < \varepsilon < 1/(2d)$ . On this random set of points  $P_1, \ldots, P_{n^{1+\varepsilon}}$  we form a random (d+1)-uniform hypergraph  $\mathcal{G}(\beta) = (V, \mathcal{E})$  with the vertices being the  $n^{1+\varepsilon}$  random points in  $[0,1]^d$ , thus  $|V| = n^{1+\varepsilon}$ . Every (d+1) vertices  $P_{i_1}, \ldots, P_{i_{d+1}}$ , of these  $n^{1+\varepsilon}$  vertices form an edge in  $\mathcal{G}(\beta)$  if the volume  $\operatorname{vol}(P_{i_1}, \ldots, P_{i_{d+1}})$  of the corresponding simplex is at most  $\beta$ , i.e.,  $\{P_{i_1}, \ldots, P_{i_{d+1}}\} \in \mathcal{E}$  if and only if  $\operatorname{vol}(P_{i_1}, \ldots, P_{i_{d+1}}) \le \beta$ . We will show for the choice  $\beta := c \cdot \log n/n^d$ , where c > 0 is a suitable constant, that among these  $n^{1+\varepsilon}$  vertices there exists an independent set of n vertices. Then, every simplex determined by (d+1) distinct points of these n points has volume at least  $\Omega(\log n/n^d)$ .

First we estimate the expected number  $E(|\mathcal{E}|)$  of edges in the random hypergraph  $\mathcal{G}(\beta)$ .

**Lemma 2.3.** For some constant  $C_d > 0$ , the expected number  $E(|\mathcal{E}|)$  of edges in the random hypergraph  $\mathcal{G}(\beta) = (V, \mathcal{E})$  satisfies:

(1) 
$$E(|\mathcal{E}|) \le C_d \cdot \beta \cdot n^{(1+\varepsilon)(d+1)}.$$

**Proof.** Our arguments are similar to those in [3]. We give an upper bound on the probability  $\operatorname{Prob}(\operatorname{vol}(P_1,\ldots,P_{d+1}) \leq \beta)$  that (d+1) points  $P_1,\ldots,P_{d+1}$  dropped in  $[0,1]^d$ , uniformly at random and independently of each other, form a simplex of volume at most  $\beta$ , i.e., we will show for some constant  $C'_d > 0$  and for every  $\beta > 0$ :

(2) 
$$\operatorname{Prob}(\operatorname{vol}(P_1,\ldots,P_{d+1}) \leq \beta) \leq C'_d \cdot \beta.$$

For  $k=2,\ldots,d+1$ , let  $x_k$  denote the Euclidean distance of  $P_k$  from the (k-2)-dimensional hyperplane  $H_{k-1}$  determined by the points  $P_1,\ldots,P_{k-1}$ . Assume that the points  $P_1,\ldots,P_{k-1},\ k=2,\ldots,d$ , are already fixed. We estimate the probability that the Euclidean distance  $x_k$  lies in the infinitesimal range  $[g_k,g_k+dg_k]$ . Taking the differences of the corresponding volumes of the cylinders determined by all points with Euclidean distance at most  $(g_k+dg_k)$  and  $g_k$ , respectively, (which are given by the volumes of (d+2-k)-dimensional balls with radii  $(g_k+dg_k)$  and  $g_k$ , respectively, multiplied by some positive

constant, which depends on d only) from the hyperplane  $H_{k-1}$ , we infer for some constant  $c_d > 0$ :

$$Prob(g_k \le x_k \le g_k + dg_k) \le d(c_d \cdot g_k^{d+2-k}) = c_d \cdot (d+2-k) \cdot g_k^{d+1-k} dg_k.$$

Now, having fixed the points  $P_1, \ldots, P_d$ , the point  $P_{d+1}$  must fulfill  $\operatorname{vol}(P_1, \ldots, P_{d+1}) \leq \beta$ , hence the Euclidean distance  $x_{d+1}$  of  $P_{d+1}$  from the hyperplane determined by  $P_1, \ldots, P_d$  must satisfy

$$\frac{1}{d!} \cdot x_{d+1} \cdot \prod_{k=2}^{d} g_k \le \beta .$$

The Euclidean distance between two points in  $[0,1]^d$  is at most  $\sqrt{d}$ , thus, the point  $P_{d+1}$  must lie within a box of base area at most  $(\sqrt{d})^{d-1}$  and of height at most

$$2 \cdot d! \cdot \frac{\beta}{\prod_{k=2}^d g_k}$$
,

which happens with probability at most

$$2 \cdot d! \cdot (\sqrt{d})^{d-1} \cdot \frac{\beta}{\prod_{k=2}^d g_k} .$$

The distances  $x_2, ..., x_d$  can be arbitrary within the range  $[0, \sqrt{d}]$ . Collecting constant factors, which only depend on the dimension d to constants  $C'_d, C''_d > 0$ , we infer

$$\operatorname{Prob}\left(\operatorname{vol}\left(P_{1},\ldots,P_{d+1}\right) \leq \beta\right)$$

$$\leq \int_{0}^{\sqrt{d}} \ldots \int_{0}^{\sqrt{d}} \left(\prod_{k=2}^{d} c_{d} \cdot (d+2-k) \cdot g_{k}^{d+1-k}\right)$$

$$\cdot \frac{2 \cdot d! \cdot (\sqrt{d})^{d-1} \cdot \beta}{\prod_{k=2}^{d} g_{k}} dg_{d} \ldots dg_{2}$$

$$= C''_{d} \cdot \beta \cdot \int_{0}^{\sqrt{d}} \ldots \int_{0}^{\sqrt{d}} \prod_{k=2}^{d} g_{k}^{d-k} dg_{d} \ldots dg_{2}$$

$$= C''_{d} \cdot \beta \cdot \frac{1}{(d-1)!} \cdot d^{\frac{d \cdot (d-1)}{4}}$$

$$= C'_{d} \cdot \beta .$$

There are  $\binom{n^{1+\varepsilon}}{d+1}$  possibilities to choose (d+1) out of the  $n^{1+\varepsilon}$  random points, hence with (2) for some constant  $C_d > 0$  the expected number  $E(|\mathcal{E}|)$ 

of edges in the random hypergraph  $\mathcal{G}(\beta) = (V, \mathcal{E})$  satisfies:

$$E(|\mathcal{E}|) \le C'_d \cdot \beta \cdot \binom{n^{1+\varepsilon}}{d+1} \le C_d \cdot \beta \cdot n^{(1+\varepsilon)(d+1)}$$
.

To apply Theorem 2.2, we will show that the expected number of 'bad configurations' among the  $n^{1+\varepsilon}$  random points is small, i.e., much less than  $n^{1+\varepsilon}$ . These bad configurations are pairs of points with small Euclidean distance and 2-cycles in the hypergraph  $\mathcal{G}(\beta)$ .

First we give an upper bound on the probability that there exist two distinct points P,Q among the  $n^{1+\varepsilon}$  random points which have Euclidean distance dist(P,Q) less than some value D>0.

**Lemma 2.4.** For every real number D > 0 and random points  $P_1, \ldots, P_{n^{1+\varepsilon}} \in [0,1]^d$  it is

(3) 
$$\operatorname{Prob}(\exists k \neq l : \operatorname{dist}(P_k, P_l) < D) \leq c_d \cdot D^d \cdot n^{2+2\varepsilon}.$$

**Proof.** For a fixed point  $P_k$ , the probability that the point  $P_l$ ,  $l \neq k$ , has Euclidean distance less than D from  $P_k$ , is given by the volume of the d-dimensional ball with center  $P_k$  and radius D, i.e., by  $c'_d \cdot D^d$  for some constant  $c'_d > 0$ . Since there are  $\binom{n^{1+\varepsilon}}{2}$  choices for the points  $P_k$  and  $P_l$ , we have for some constant  $c_d > 0$ :

$$\operatorname{Prob}(\exists k \neq l : \operatorname{dist}(P_k, P_l) < D)$$

$$\leq \sum_{1 \leq k < l \leq n^{1+\varepsilon}} \operatorname{Prob}(\operatorname{dist}(P_k, P_l) < D)$$

$$\leq \binom{n^{1+\varepsilon}}{2} \cdot c'_d \cdot D^d \leq c_d \cdot D^d \cdot n^{2+2\varepsilon}.$$

With (3) and  $D_0 := n^{-2/(d-1)}$ , where  $0 < \varepsilon < 2/(d-1)$ , we obtain that

$$Prob(\exists k \neq l : dist(P_k, P_l) < D_0) = o(1),$$

thus,

(4) 
$$\operatorname{Prob}(\forall k \neq l : \operatorname{dist}(P_k, P_l) \geq D_0) = 1 - o(1)$$
,

and with probability close to 1 distinct points have Euclidean distance at least  $D_0$ .

Next, for j = 2, ..., d, we will give an upper bound on the conditional expected numbers

$$E(s_{2,j}(\mathcal{G}(\beta)) \mid \forall k \neq l : \operatorname{dist}(P_k, P_l) \geq D_0)$$

of (2, j)-cycles in  $\mathcal{G}(\beta)$ , that is, the expected numbers of pairs  $\{E_1, E_2\}$  of edges  $E_1, E_2 \in \mathcal{E}$  with  $|E_1 \cap E_2| = j$ , given that distinct points have Euclidean distance at least  $D_0$ .

**Lemma 2.5.** For j = 2, ..., d-1 and constants  $c_j(d) > 0$  the random hypergraph  $\mathcal{G}(\beta)$  satisfies:

(5) 
$$E(s_{2,j}(\mathcal{G}(\beta)) | \forall k \neq l : \operatorname{dist}(P_k, P_l) \geq D_0) \leq c_j(d) \cdot \beta^2 \cdot n^{(1+\varepsilon)(2d+2-j)}$$
, and for  $j = d$  and a constant  $c(d) > 0$  it is

(6) 
$$E(s_{2,d}(\mathcal{G}(\beta)) \mid \forall k \neq l : \operatorname{dist}(P_k, P_l) \geq D_0) \leq c(d) \cdot \beta^2 \cdot n^{(1+\varepsilon)(d+2)} \cdot \log n$$
.

**Proof.** Let  $j=2,\ldots,d$ . Consider (2d+2-j) random points  $P_1,\ldots,P_{2d+2-j}\in [0,1]^d$  where the Euclidean distances satisfy  $\operatorname{dist}(P_k,P_l)\geq D_0=n^{-2/(d-1)}$  for  $1\leq k< l\leq 2d+2-j$ . We will give an upper bound on the following conditional probability:

$$\operatorname{Prob}(P_1,\ldots,P_{2d+2-j})$$
 form a  $(2,j)$ -cycle in  $\mathcal{G}(\beta) \mid \forall k \neq l : \operatorname{dist}(P_k,P_l) \geq D_0$ .

Let us assume that the two simplices, which yield a (2,j)-cycle, are  $E = \{P_1, \dots, P_{d+1}\} \in \mathcal{E}$  and  $E' = \{P_1, \dots, P_j, P_{d+2}, P_{d+3}, \dots, P_{2d+2-j}\} \in \mathcal{E}$  with

$$(7) vol(P_1, \dots, P_{d+1}) \le \beta$$

and

(8) 
$$\operatorname{vol}(P_1, \dots, P_j, P_{d+2}, \dots, P_{2d+2-j}) \leq \beta$$
.

All possibilities for forming a (2,j)-cycle will be taken into account by the constant factor  $\binom{2d+2-j}{d+1} \cdot \binom{d+1}{j}$ . Let  $\mathcal{F}_{E,E'}$  denote the event " $\{E,E'\}$  is a (2,j)-cycle in  $\mathcal{G}(\beta)$  given that  $\forall k \neq l$ : dist $(P_k,P_l) \geq D_0$ ". We will estimate the probability  $\text{Prob}(\mathcal{F}_{E,E'})$ .

For  $k=2,\ldots,d+1$ , let  $x_k$  denote the Euclidean distance of the point  $P_k$  from the hyperplane determined by  $P_1,\ldots,P_{k-1}$ . For  $l=d+2,\ldots,2d+2-j$ , let  $y_l$  be the Euclidean distance of the point  $P_l$  from the hyperplane determined by  $P_1,\ldots,P_j,P_{d+2},\ldots,P_{l-1}$ , where for l=d+2 the hyperplane is determined by  $P_1,\ldots,P_j$ . Assume that the points  $P_1,\ldots,P_{k-1}$ , are already fixed. As in the proof of Lemma 2.3 we have for some constant  $c'_d>0$ :

$$\operatorname{Prob}(g_k \le x_k \le g_k + dg_k) \le d(c_d \cdot g_k^{d+2-k}) \le c'_d \cdot g_k^{d+1-k} dg_k.$$

Also, for  $l = d+2, \ldots, 2d+1-j$ , given the points  $P_1, \ldots, P_j, P_{d+2}, \ldots, P_{l-1}$  we have

$$Prob(h_l \le y_l \le h_l + dh_l) \le d(c_d \cdot h_l^{2d-l-j+3}) \le c'_d \cdot h_l^{2d-l-j+2} dh_l.$$

To satisfy (7), given the points  $P_1, P_2, \ldots, P_d$ , the point  $P_{d+1}$  must lie in a box of volume at most

$$C'_d \cdot \frac{\beta}{\prod_{k=2}^d g_k}$$
,

where  $C'_d>0$  is a constant. Similarly, if the points  $P_1, \ldots, P_j, P_{d+2}, \ldots, P_{2d+1-j}$  are already fixed, to satisfy (8), the point  $P_{2d+2-j}$  must lie in a box of volume at most

$$C'_d \cdot \frac{\beta}{\prod_{k=2}^{j} g_k \cdot \prod_{l=d+2}^{2d+1-j} h_l}$$
.

We infer for some constant  $C_d > 0$ :

$$\begin{aligned}
& \text{Prob}(\mathcal{F}_{E,E'}) \\
&= \text{Prob}\left(\{E,E'\} \text{ is a } (2,j)\text{-cycle in } \mathcal{G}(\beta) \mid \forall k \neq l : \text{dist}(P_k,P_l) \geq D_0\right) \\
&\leq C_d \cdot \int_{D_0}^{\sqrt{d}} \cdots \int_{D_0}^{\sqrt{d}} \left(\prod_{k=2}^d g_k^{d+1-k}\right) \cdot \left(\prod_{l=d+2}^{2d+1-j} h_l^{2d-l-j+2}\right) \cdot \frac{\beta^2}{(\prod_{k=2}^d g_k) \cdot (\prod_{k=2}^j g_k) \cdot (\prod_{l=d+2}^{2d+1-j} h_l)} \cdot dh_{2d+1-j} \cdots dh_{d+2} \, dg_d \cdots dg_2 \\
&= C_d \cdot \beta^2 \cdot \int_{D_0}^{\sqrt{d}} \cdots \int_{D_0}^{\sqrt{d}} \left(\prod_{k=2}^j g_k^{d-1-k}\right) \cdot \left(\prod_{k=j+1}^d g_k^{d-k}\right) \cdot \left(\prod_{l=d+2}^d h_l^{2d-l-j+1}\right) \, dh_{2d+1-j} \cdots dh_{d+2} \, dg_d \cdots dg_2 \, .
\end{aligned}$$

For nonnegative exponents the terms  $g_k^{d-1-k}$  and  $h_l^{2d-l-j+1}$  contribute with respect to the integration at most a constant factor dependent on d only. Only in the case k=j=d the exponent (d-1-k) of  $g_k=g_d$  is negative. Hence, for  $j=2,\ldots,d-1$ , we have for some constant  $C_d^*>0$ 

(9) 
$$\operatorname{Prob}(\mathcal{F}_{E,E'}) \le C_d^* \cdot \beta^2,$$

while for j = d, and here we use the assumption  $D_0 = n^{-2/(d-1)}$ , we obtain for some constants  $C'_d, C''_d, C^{**}_d > 0$ :

(10) 
$$\operatorname{Prob}(\mathcal{F}_{E,E'}) \leq C'_d \cdot \beta^2 \cdot \int_{D_0}^{\sqrt{d}} \frac{1}{g_d} dg_d \leq C''_d \cdot \beta^2 \cdot \log(1/D_0)$$
$$\leq C_d^{**} \cdot \beta^2 \cdot \log n.$$

We can choose (2d+2-j) points from  $n^{1+\varepsilon}$  points in  $\binom{n^{1+\varepsilon}}{2d+2-j}$  ways. Taking into account the number  $\binom{2d-j+2}{d+1} \cdot \binom{d+1}{j}$  of possibilities to form a (2,j)-cycle, we conclude with (9) for  $j=2,\ldots,d-1$ , that the conditional expected numbers  $E(s_{2,j}(\mathcal{G}(\beta)) \mid \forall k \neq l : \operatorname{dist}(P_k,P_l) \geq D_0)$  of (2,j)-cycles in  $\mathcal{G}(\beta)$  satisfy for constants  $c_j(d),c(d)>0$ :

$$E(s_{2,j}(\mathcal{G}(\beta)) \mid \forall k \neq l : \operatorname{dist}(P_k, P_l) \geq D_0)$$

$$\leq \binom{2d-j+2}{d+1} \cdot \binom{d+1}{j} \cdot C_d^* \cdot \beta^2 \cdot \binom{n^{1+\varepsilon}}{2d+2-j}$$

$$\leq c_j(d) \cdot \beta^2 \cdot n^{(1+\varepsilon)(2d+2-j)},$$

and for j = d we have by (10) that

$$E(s_{2,d}(\mathcal{G}(\beta)) \mid \forall k \neq l : \operatorname{dist}(P_k, P_l) \geq D_0)$$

$$\leq \binom{2d - j + 2}{d + 1} \cdot \binom{d + 1}{j} \cdot C_d^{**} \cdot \beta^2 \cdot \log n \cdot \binom{n^{1+\varepsilon}}{d + 2}$$

$$\leq c(d) \cdot \beta^2 \cdot n^{(1+\varepsilon)(d+2)} \cdot \log n.$$

Now we set

$$\beta := \frac{\log n}{n^d} \ .$$

**Lemma 2.6.** For fixed  $\varepsilon$  with  $0 < \varepsilon < 1/(2d)$ , there exists a hypergraph  $\mathcal{G}(\beta) = (V, \mathcal{E})$  which satisfies:

$$|V| = n^{1+\varepsilon}$$

$$|\mathcal{E}| \le 2 \cdot C_d \cdot \beta \cdot n^{(1+\varepsilon)(d+1)}$$

$$s_{2,j}(\mathcal{G}(\beta)) \le n \quad \text{for } j = 2, \dots, d.$$

**Proof.** We will show that the event  $\mathcal{F} = (|\mathcal{E}| \leq 2 \cdot C_d \cdot \beta \cdot n^{(1+\varepsilon)(d+1)})$  and  $(\forall k \neq l : \text{dist}(P_k, P_l) \geq D_0)$  and  $(\forall j : s_{2,j}(\mathcal{G}(\beta)) \leq n)$ " happens with positive probability for our random hypergraph  $\mathcal{G}(\beta) = (V, \mathcal{E})$ .

The complementary event of  $\mathcal{F}$  is  $\overline{\mathcal{F}} = \text{``}(|\mathcal{E}| > 2 \cdot C_d \cdot \beta \cdot n^{(1+\varepsilon)(d+1)})$  or  $(\exists k \neq l : \text{dist}(P_k, P_l) < D_0)$  or  $(\exists j : [s_{2,j}(\mathcal{G}(\beta)) > n \text{ and } \forall k \neq l : \text{dist}(P_k, P_l) \geq D_0])$ ".

Using (1), (4) and Markov's inequality, i.e.,  $\operatorname{Prob}(X \geq \alpha) \leq E(X)/\alpha$  for every real  $\alpha > 0$  and every nonnegative random variable X, we infer

$$\operatorname{Prob}(\overline{\mathcal{F}}) \leq \operatorname{Prob}(|\mathcal{E}| > 2 \cdot C_d \cdot \beta \cdot n^{(1+\varepsilon)(d+1)}) + \operatorname{Prob}(\exists k \neq l : \operatorname{dist}(P_k, P_l) < D_0) +$$

$$+\operatorname{Prob}\left(\exists j: [s_{2,j}(\mathcal{G}(\beta)) > n \text{ and } \forall k \neq l: \operatorname{dist}(P_k, P_l) \geq D_0]\right)$$

$$\leq \frac{1}{2} + o(1) + \sum_{j=2}^{d} \operatorname{Prob}\left(s_{2,j}(\mathcal{G}(\beta)) > n \text{ and } \forall k \neq l: \operatorname{dist}(P_k, P_l) \geq D_0\right)$$

$$= \frac{1}{2} + o(1) + \sum_{j=2}^{d} \operatorname{Prob}\left(\forall k \neq l: \operatorname{dist}(P_k, P_l) \geq D_0\right) \cdot \cdot \cdot \operatorname{Prob}\left(s_{2,j}(\mathcal{G}(\beta)) > n \mid \forall k \neq l: \operatorname{dist}(P_k, P_l) \geq D_0\right)$$

$$= \frac{1}{2} + o(1) + (1 - o(1)) \cdot \cdot \cdot \cdot \sum_{j=2}^{d} \operatorname{Prob}\left(s_{2,j}(\mathcal{G}(\beta)) > n \mid \forall k \neq l: \operatorname{dist}(P_k, P_l) \geq D_0\right)$$

$$(12) \leq \frac{1}{2} + o(1) + (1 - o(1)) \cdot \cdot \cdot \sum_{j=2}^{d} \frac{E(s_{2,j}(\mathcal{G}(\beta))) \mid \forall k \neq l: \operatorname{dist}(P_k, P_l) \geq D_0}{n}.$$

For  $j=2,\ldots,d-1$  we have by (5) and (11) for  $\varepsilon < 1/(2d)$ :

$$\frac{E(s_{2,j}(\mathcal{G}(\beta)) \mid \forall k \neq l : \operatorname{dist}(P_k, P_l) \geq D_0)}{n}$$

$$\leq \frac{c_j(d) \cdot \beta^2 \cdot n^{(1+\varepsilon)(2d+2-j)}}{n}$$

$$= c_j(d) \cdot (\log n)^2 \cdot n^{1-j+\varepsilon(2d+2-j)}$$

$$= o(1).$$

and for j=d and  $\varepsilon < (d-1)/(d+2)$  we have by (6) and (11):

$$\frac{E(s_{2,j}(\mathcal{G}(\beta)) \mid \forall k \neq l : \operatorname{dist}(P_k, P_l) \geq D_0)}{n}$$

$$\leq \frac{c(d) \cdot \beta^2 \cdot n^{(1+\varepsilon)(d+2)} \cdot \log n}{n}$$

$$= c(d) \cdot (\log n)^3 \cdot n^{-d+1+\varepsilon(d+2)}$$

$$= o(1) .$$

We conclude with (12) that  $\operatorname{Prob}(\overline{\mathcal{F}}) \leq 1/2 + o(1)$  and hence  $\operatorname{Prob}(\mathcal{F}) > 0$  for  $0 < \varepsilon < 1/(2d)$ . Thus there exists a desired hypergraph  $\mathcal{G}(\beta) = (V, \mathcal{E})$ .

We take the (d+1)-uniform hypergraph  $\mathcal{G}(\beta)$  with  $0 < \varepsilon < 1/(2d)$ , which exists by Lemma 2.6, and we remove one vertex from each (2,j)-cycle,  $j = 2,3,\ldots,d$ . We obtain an induced subhypergraph  $\mathcal{G}_1(\beta) = (V_1,\mathcal{E}_1)$  of  $\mathcal{G}(\beta) = (V,\mathcal{E})$  with  $|V_1| = (1-o(1)) \cdot n^{1+\varepsilon}$  vertices and  $|\mathcal{E}_1| \leq 2 \cdot C_d \cdot \beta \cdot n^{(1+\varepsilon)(d+1)}$ 

and without any 2-cycles. Hence,  $\mathcal{G}_1(\beta)$  has average degree at most  $t^d = 2 \cdot C_d \cdot (1+o(1)) \cdot (d+1) \cdot \beta \cdot n^{(1+\varepsilon)d}$ . Set  $c_d^{**} := (2 \cdot C_d \cdot (1+o(1)) \cdot (d+1))^{1/d}$ .

We apply Theorem 2.2 to the (d+1)-uniform subhypergraph  $\mathcal{G}_1(\beta) = (V_1, \mathcal{E}_1)$  and by the choice of  $\beta$  in (11) the independence number  $\alpha(\mathcal{G}_1(\beta))$  satisfies for suitable constants  $c'_d, c^*_d > 0$ :

$$\alpha(\mathcal{G}(\beta)) \ge \alpha(\mathcal{G}_1(\beta)) \ge c_{d+1} \cdot \frac{(1 - o(1)) \cdot n^{1+\varepsilon}}{c_d^{**} \cdot \beta^{1/d} \cdot n^{1+\varepsilon}} \cdot \left(\log(c_d^{**} \cdot \beta^{1/d} \cdot n^{1+\varepsilon})\right)^{1/d}$$

$$\ge c_d' \cdot \frac{n}{(\log n)^{1/d}} \cdot (\log n^{\varepsilon})^{1/d}$$

$$\ge c_d^* \cdot \frac{n}{(\log n)^{1/d}} \cdot (\log n)^{1/d}$$

$$\ge c_d^* \cdot n.$$

Thus, among the  $n^{1+\varepsilon}$  points in  $[0,1]^d$  there is a subset of  $c_d^* \cdot n$  points, such that each simplex spanned by any (d+1) of these  $c_d^* \cdot n$  points has volume at least  $\beta = \log n/n^d$ . By adapting constant factors, i.e., choosing  $\beta = c \cdot \log n/n^d$  for a suitable constant c > 0, there exist n points in  $[0,1]^d$  such that the volume of every simplex spanned by any (d+1) of these n points is at least  $\Omega(\log n/n^d)$ . This finishes the proof of Theorem 1.1.

## 3. Concluding Remarks

We showed by a probabilistic argument the existence of a configuration of n points in the d-dimensional unit cube  $[0,1]^d$  such that the volume of every simplex formed by any (d+1) of these points is at least  $\Omega(\log n/n^d)$ . Although there is an algorithmic version of Theorem 2.2 available, see [8] and [4], it seems to be difficult and involved, to turn our arguments into a deterministic polynomial time algorithm. For the 2-dimensional case we succeeded in doing so by using a sufficiently fine grid [5], and very recently also for the case d=3, see [11]. Moreover, it would also be interesting to investigate upper bounds for the d-dimensional version of Heilbronn's problem.

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